

How far can the Rödl nibble go?

Tom Kelly

Joint work with Stephen Gould



Georgia Institute
of Technology



Combinatorics at the Confluence
Carnegie Mellon University, University of Pittsburgh
July 20, 2026

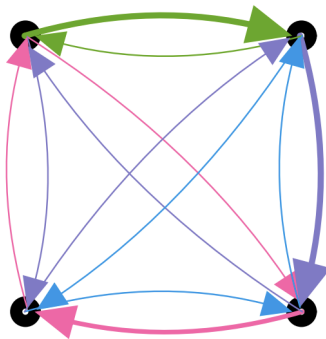
Part I

Applications

Let G be a directed graph.

proper arc-coloring: at each vertex, all outgoing arcs receive distinct colors, and all incoming arcs receive distinct colors

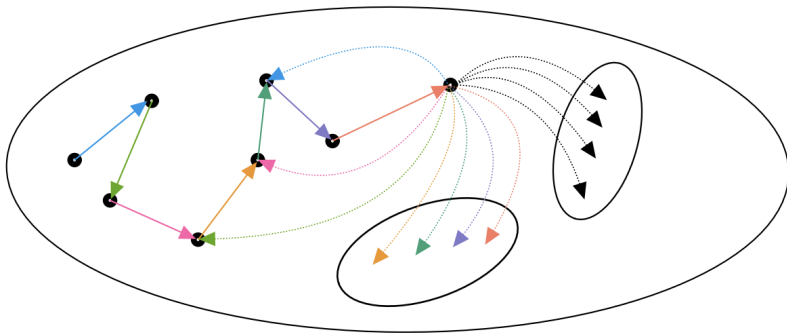
rainbow: edges have distinct colors



A rainbow path of length three in a proper arc-coloring of $K_4^{\leftrightarrow}$.

Question

Given a proper n -arc-coloring of $K_n^{\leftrightarrow}$, what is the expected length of a rainbow self-avoiding walk constructed by random greedy (as $n \rightarrow \infty$)?



Question

Given a proper n -arc-coloring of $K_n^{\leftrightarrow}$, what is the expected length of a rainbow self-avoiding walk constructed by random greedy (as $n \rightarrow \infty$)?

Heuristics suggest: $n - \sqrt{n}$

Let $p = (1 - t/n)$. At time t , pn vertices and colors remain unused, and they should “look like” a set chosen independently with probability p .

In the independent setting,

$$\mathbb{E} [\# \text{ valid choices}] = p^2 n,$$

and

$$p^2 n < 1 \quad \Rightarrow \quad p < 1/\sqrt{n}.$$

Conjecture (Gyárfás & Sárközy, 2014;, Gould & Kelly, 2023)

Every properly arc-colored $K_n^{\leftrightarrow}$ contains a rainbow path or cycle of length at least $n - 1$.

- If true, implies both **Andersen's conjecture (2005)** and the even- n case of the **Ryser–Brualdi–Stein conjecture (1967)**.

Benzing, Pokrovskiy, and Sudakov (2020): $n - O(n^{4/5})$

Balogh and Molla (2019): $n - O(\log n \sqrt{n})$ for undirected graphs

Gould and Kelly (2023): true for *almost all* proper n -arc-colorings

Conjecture (Gyárfás & Sárközy, 2014;, Gould & Kelly, 2023)

Every properly arc-colored $K_n^{\leftrightarrow}$ contains a rainbow path or cycle of length at least $n - 1$.

- If true, implies both **Andersen's conjecture (2005)** and the even- n case of the **Ryser–Brualdi–Stein conjecture (1967)**.

Benzing, Pokrovskiy, and Sudakov (2020): $n - O(n^{4/5})$

Balogh and Molla (2019): $n - O(\log n \sqrt{n})$ for undirected graphs

Gould and Kelly (2023): true for *almost all* proper n -arc-colorings

Theorem (Gould and Kelly, 2025+)

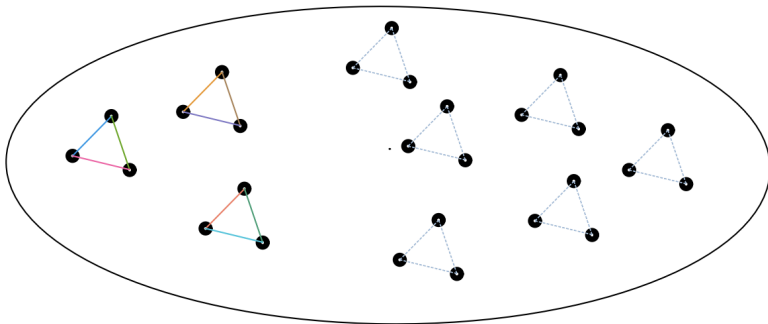
Every properly n -arc-colored $K_n^{\leftrightarrow}$ contains a rainbow cycle of length at least $n - n^{1/2+o(1)}$.

- probabilistic proof – matches heuristic

Conjecture (Müeyesser, 2023)

Every properly n -edge-colored K_n contains a rainbow triangle tiling covering all but $O(1)$ vertices.

- Motivated by **Friedlander-Gordon-Tannenbaum conjecture (1981)** from Group Theory.



Conjecture (Müyesser, 2023)

Every properly n -edge-colored K_n contains a rainbow triangle tiling covering all but $O(1)$ vertices.

Question

What is the expected number of vertices left uncovered by a rainbow triangle tiling constructed by random greedy (as $n \rightarrow \infty$)?

Heuristics suggest: $n^{3/5}$

If vtc's and colors are chosen independently with prob. p , then $\forall$ vertex v ,

$$\mathbb{E}[\# \text{ triangles containing } v] = p^5 \binom{n}{2},$$

and

$$p^5 n^2 < 1 \quad \Rightarrow \quad p < 1/n^{2/5}.$$

Conjecture (Müyesser, 2023)

Every properly n -edge-colored K_n contains a rainbow triangle tiling covering all but $O(1)$ vertices.

Theorem (Gould and Kelly, 2025+)

Every properly n -edge-colored K_n has a rainbow triangle tiling covering all but at most $n^{3/5+o(1)}$ vertices.

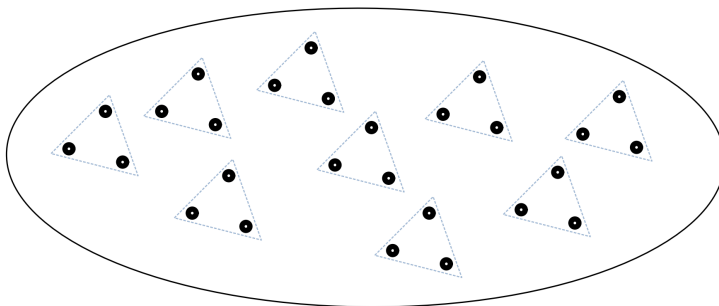
- The best that can be obtained with prior results is $n^{3/4+o(1)}$.

Part II

Hypergraph matchings

matching: set of pairwise disjoint edges

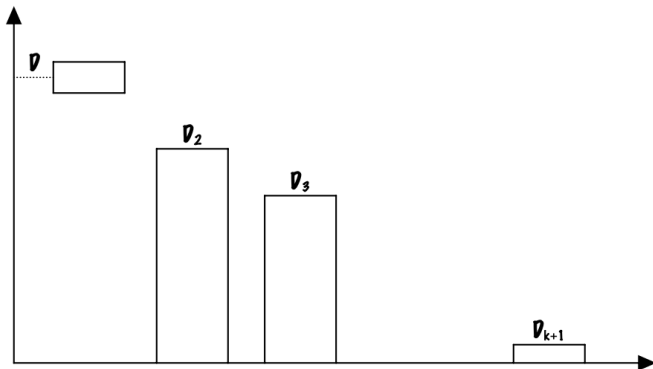
k -uniform: every edge has size k



A 3-uniform matching

In what follows, H is an n -vertex $(k + 1)$ -uniform hypergraph which is

- $(1 \pm \varepsilon)D$ -regular, and
- every set of i vertices is contained in at most D_i edges of H .



Pippenger's Theorem (Füredi, 1988)

If $\varepsilon \ll \gamma$ and $D_2 \leq \varepsilon D$, then H has a matching covering all but $\leq \gamma n$ vtc's.

Heuristic: $D_2 \ll D \Rightarrow$ when pn vertices remain in nibble/random greedy, vertices “look like” they are chosen independently with probability p . In the independent setting,

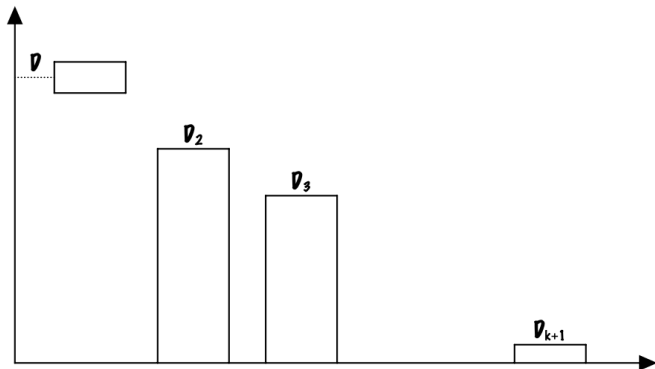
$$\mathbb{E}[\text{degree of } v] = p^k D,$$

and

$$p^k D \gg D_2 \quad \Rightarrow \quad p \gg (D_2/D)^{1/k}$$

Theorem (Kostochka and Rödl, 1997)

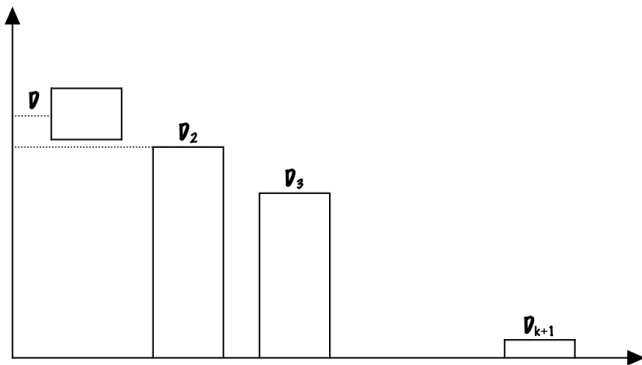
If $\varepsilon \leq O(\sqrt{\log D/D})$ and $D_2 \leq D^{1-o(1)}$, then H has a matching covering all but $\leq (D_2/D)^{1/k+o(1)} n$ vtcs.



The codegree sequence before nibble

Theorem (Kostochka and Rödl, 1997)

If $\varepsilon \leq O(\sqrt{\log D/D})$ and $D_2 \leq D^{1-o(1)}$, then H has a matching covering all but $\leq (D_2/D)^{1/k+o(1)} n$ vtcs.



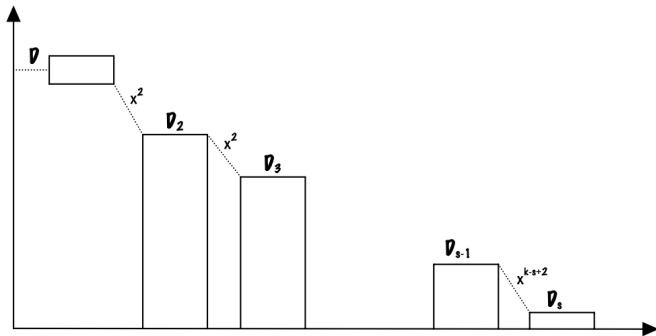
The codegree sequence after nibble

Theorem (Vu, 2000)

If there exists $x > 0$ and $s \in [k + 1]$ such that

- **(well-spaced codegrees)** $x^2 \leq D_j/D_{j+1}$ for all $j \in [s - 1]$;
- **(large penultimate codegree gap)** $x^{k-s+2} \leq D_{s-1}/D_s$;
- **(nearly regular)** $\varepsilon \leq 1/x$,

then H has a matching covering all but $\leq n(\log^{O(1)} D)/x$ vtc's.



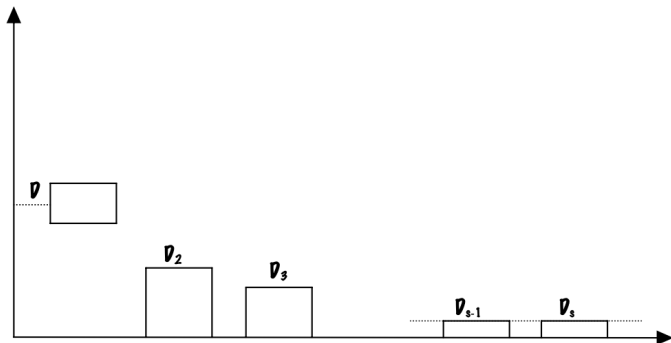
The codegree sequence before nibble

Theorem (Vu, 2000)

If there exists $x > 0$ and $s \in [k + 1]$ such that

- **(well-spaced codegrees)** $x^2 \leq D_j/D_{j+1}$ for all $j \in [s - 1]$;
- **(large penultimate codegree gap)** $x^{k-s+2} \leq D_{s-1}/D_s$;
- **(nearly regular)** $\varepsilon \leq 1/x$,

then H has a matching covering all but $\leq n(\log^{O(1)} D)/x$ vtc's.



The codegree sequence after nibble

Theorem (Vu, 2000)

If there exists $x > 0$ and $s \in [k + 1]$ such that

- **(well-spaced codegrees)** $x^2 \leq D_j/D_{j+1}$ for all $j \in [s - 1]$;
- **(large penultimate codegree gap)** $x^{k-s+2} \leq D_{s-1}/D_s$;
- **(nearly regular)** $\varepsilon \leq 1/x$,

then H has a matching covering all but $\leq n(\log^{O(1)} D)/x$ vtcs.

Our idea: Can we iterate? Reapply Vu with $s \leftarrow s - 1$?

Theorem (Vu, 2000)

If there exists $x > 0$ and $s \in [k + 1]$ such that

- **(well-spaced codegrees)** $x^2 \leq D_j/D_{j+1}$ for all $j \in [s - 1]$;
- **(large penultimate codegree gap)** $x^{k-s+2} \leq D_{s-1}/D_s$;
- **(nearly regular)** $\varepsilon \leq 1/x$,

then H has a matching covering all but $\leq n(\log^{O(1)} D)/x$ vtcs.

Our idea: Can we iterate? Reapply Vu with $s \leftarrow s - 1$?

Two main issues:

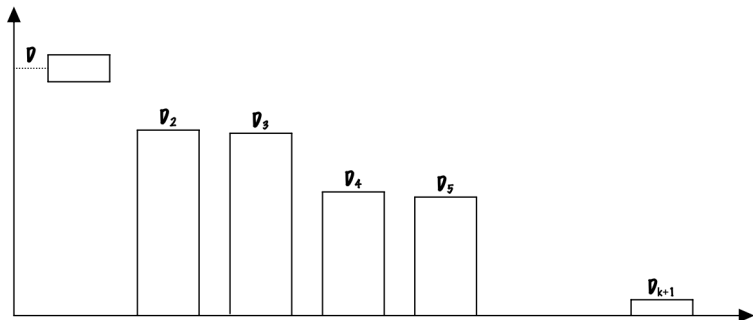
- The remaining hypergraph is not sufficiently regular.
- Vu's result cannot handle “clustered codegrees”.

Theorem (Gould and Kelly, 2025+)

If there exists $B > 0$ such that

- (large degree/codegree ratios) $B \leq (D/D_j)^{\frac{1}{j-1}} \quad \forall j \geq 4$;
- (large 2-degree gap) $B \leq \sqrt{D/D_2}$;
- (nearly regular) $\varepsilon \leq 1/B$,

then H has a matching covering all but $\leq n(\log^{O(1)} D)/B^{1-o(1)}$ vtc's.



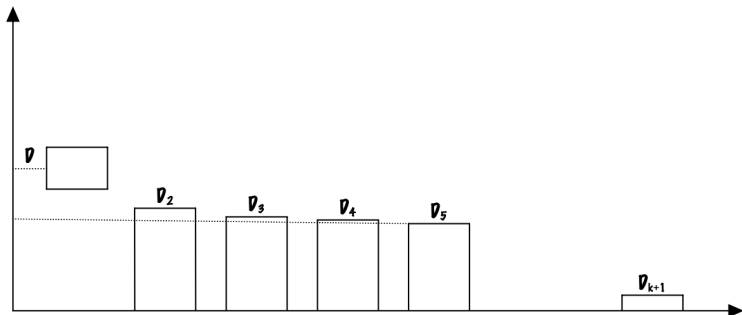
The codegree sequence before nibble

Theorem (Gould and Kelly, 2025+)

If there exists $B > 0$ such that

- (large degree/codegree ratios) $B \leq (D/D_j)^{\frac{1}{j-1}} \quad \forall j \geq 4$;
- (large 2-degree gap) $B \leq \sqrt{D/D_2}$;
- (nearly regular) $\varepsilon \leq 1/B$,

then H has a matching covering all but $\leq n(\log^{O(1)} D)/B^{1-o(1)}$ vtc's.



The codegree sequence after nibble

Theorem (Gould and Kelly, 2025+)

If there exists $B > 0$ such that

- (large degree/codegree ratios) $B \leq (D/D_j)^{\frac{1}{j-1}} \quad \forall j \geq 4$;
- (large 2-degree gap) $B \leq \sqrt{D/D_2}$;
- (nearly regular) $\varepsilon \leq 1/B$,

then H has a matching covering all but $\leq n(\log^{O(1)} D)/B^{1-o(1)}$ vtc's.

Heuristic: $D_{j+1} \ll D_j \Rightarrow$ when pn vertices remain in nibble/random greedy:

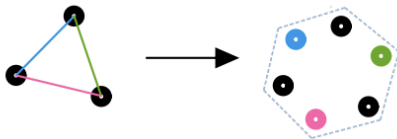
$$\mathbb{E}[\text{degree of a } j\text{-set}] \leq p^{k+1-j} D_j,$$

and

$$p^{k+1-j} D_j \ll p^k D \quad \Rightarrow \quad p \gg (D_j/D)^{\frac{1}{j-1}}$$

Theorem (Gould and Kelly, 2025+)

Every properly n -edge-colored K_n has a rainbow triangle tiling covering all but at most $n^{3/5+o(1)}$ vertices.



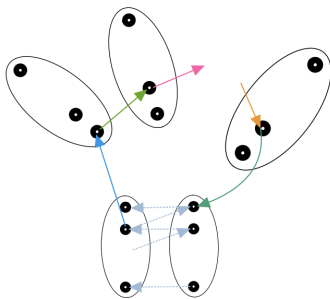
Rainbow triangle tilings $\longleftrightarrow$ matchings in H ,

where H is 6-uniform, $\binom{n}{2}$ -regular, with codegree sequence

$$\left(\binom{n}{2}, n-2, n-2, 1, 1 \right)$$

Theorem (Gould and Kelly, 2025+)

Every properly n -arc-colored $K_n^{\leftrightarrow}$ contains a rainbow cycle of length at least $n - n^{1/2+o(1)}$.



Long rainbow cycles $\longleftarrow$ “A-perfect” matchings in H ,
 where H is $2\ell + 1$ -uniform, $\approx c_\ell n^\ell$ -regular, with codegree sequence

$$(c_\ell n^\ell, n^{\ell-1}, n^{\ell-1}, \dots, n, n, 1, 1, 1).$$

Other applications of our result: Improvements to

- **Molloy–Reed (2000)** bound on chromatic index of uniform hypergraphs that are far from linear;
- **Bohman–Newman (2022)** asymptotics for maximum diameter of a d -dimensional strongly connected simplicial complex:
 - ▶ **Glock–Parczyk–Rathke–Szabó (2026+)**: tight bound;
- **Kang–Kühn–Methuku–Osthus (2023)** bound on the size of a maximum matching and the chromatic index of an (n, q, r) -Steiner system for $r > 2$.

Tiers of generality for nearly perfect hypergraph matchings:

- pseudorandom nearly perfect hypergraph matchings:
 - ▶ **Ehard–Glock–Joos (2020); Alon–Yuster (2005);**
- tripartite/reserve perfect matchings:
 - ▶ **Delcourt–Postle (2024+); Joos–Mubayi–Smith (2024+) ;**
- bipartite “A-perfect matchings”:
 - ▶ **Delcourt–Postle (2022+);**
- asymptotically good list edge-coloring:
 - ▶ **Kahn (1996);**
- asymptotically good proper edge-coloring:
 - ▶ **Pippenger–Spencer (1989);**
- nearly perfect matchings:
 - ▶ **Frankl–Rödl (1985); Pippenger (1988).**

Thanks for listening!